## FAQs on Convex Optimization

1. What is a convex programming problem?

A convex programming problem is the minimization of a convex function on a convex set, i.e.

min f(x)  
$$X \in C$$

where f:  $R^n \ \exists R$  and C  $R^n$ . f is a convex function and c a convex set. Usually C is described as follows

$$C = \{ x: g_i (x) \le 0, i=1,...,m, h_j (^)=0, j=1,...,m \}$$
  
where  $g_i's$  are convex function and  $hj's$  are affine function.

2. What is the importance of convex optimization problems?

The major importance of convex programming or convex optimization arises from the fact that every local minimum is a global minimum.



Let us consider minimizing f:  $\mathbb{R}^n \to \mathbb{R}$  or  $\mathbb{C} - \mathbb{R}^n$  where f is a convex function and C is a convex set.

Let  $\stackrel{\acute{\chi}}{}$  be a local minimum of f on C. thus  $\exists \delta > 0$  such that  $\forall z \in (\stackrel{\acute{\chi}}{}) \cap C$ ,  $f(z) \geq f(\stackrel{\acute{\chi}}{})$ . Let  $X \stackrel{\leftarrow}{} C($  take it outside  $\stackrel{B_{\delta}}{} (\stackrel{\acute{\chi}}{}) \cap C)$ . Join x &  $\stackrel{B_{\delta}}{} (\stackrel{\acute{\chi}}{}) \cap C)$ using a line segment. Let

$$Z_{\lambda} = \lambda x + (1 - \lambda) \dot{x}$$

Thus  $\exists \lambda_0 \in (0,1)$  such that  $\forall \in (0, \lambda_0, \mathcal{L}_\lambda \in B_\delta)$   $(x \to 1) \cap C$ Thus for  $\lambda \in (0, \lambda_0 \mathcal{L})$ 

$$f(\lambda x + (1-\lambda)) \xrightarrow{X} ) \ge f(X)$$

By convexity of

$$\lambda f(x) + (1-\lambda) f() \ge f(\overset{\chi}{x})$$

$$=> \lambda(f(x) - f(\overset{\chi}{x})) \ge 0$$

$$=> (f(x) - f(\overset{\chi}{x})) \ge 0, \text{ as } \lambda > 0$$

$$\overset{\chi}{x}$$

Since x is arbitrary we have as  $\uparrow$  the global minimum.

- 3. What can we tell about the continuity and differentiality of a convex function?
  - If f:  $R^n \rightarrow R$  is convex then f is continuous and even locally Lipschitz,

i.e; for any 
$$x \in \mathbb{R}^{n}$$
  
and  $K \ge 0$   
such that for all  $y, z \in B_{\delta}$  (x) we have  
 $|f(y)-f(z)| \le K ||y-z||,$ 

- If f: C  $\rightarrow$  R is convex and C is a closed convex set then, f is continuous on the interior of C.
- If f: <sup>R<sup>n</sup></sup> →R is convex, then it is differentiable almost everywhere,
   i.e.; the set of points in <sup>R<sup>n</sup></sup> at which f is not differentiable forms a set of measure zero.
- A differentiable function f:  $\mathbb{R}^n \to \mathbb{R}$  is convex if and only if; for all x, y in  $\mathbb{R}^n$ .

 $f(y) - f(x) \ge \langle \nabla f(x), y - x \rangle$ 

Thus if  $(x \in \mathbb{R}^n)$  be such that  $\nabla f = 0$ , then f has a global minimizer at x.

4. If f:  $R^n \rightarrow R$  is differentiable then can we detect it.

If f is twice continuously differentiable then there is at least a theoretical way to detect it.

- > A function f is convex if and only if the Hersian matrix  $\nabla^2$  f(x) is positive for all x  $\in \mathbb{R}^n$  semi-definitely.
- > If  $\nabla^2$  f(x) is positive definite for all x  $\in \mathbb{R}^n$ , then f is strictly converse. The converse need not be true. Example : f(x) = X<sup>4</sup>, X  $\in \mathbb{R}$
- > If f is strongly convex then  $\nabla^2$  f(x) is always positive definite.

Let f be a p-strongly convex function. since f is twice continuously differentiable, it is differentiable and hence

f(y) - f(x) 
$$\langle \nabla f(x), y - x \rangle + p$$
 ||y-x||<sup>2</sup>, P>0

Now by Taylor's theorem for any  $\lambda > 0$ , & w  $\in \mathbb{R}^n$ 

$$\mathbf{f} \stackrel{x+\lambda w}{(\mathcal{L})} = \mathbf{f} \stackrel{x}{(\mathcal{L})} + \lambda \quad \langle \nabla f(x), w \rangle + \frac{1}{2} \quad \lambda^2 \quad \langle w, \nabla \quad {}^2\mathbf{f}(x) w \quad \overset{\lambda}{\rangle} + 0 \mathcal{L} \quad {}^2\mathbf{f}(x) w \quad \overset{\lambda}{\rangle} + 0 \mathcal{L} \quad {}^2\mathbf{f}(x) = \mathbf{f} \quad \lambda^2 \quad \langle w, \nabla \quad {}^2\mathbf{f}(x) w \quad \overset{\lambda}{\rangle} + 0 \mathcal{L} \quad {}^2\mathbf{f}(x) = \mathbf{f} \quad {}^2\mathbf{f}(x) = \mathbf{f} \quad \lambda^2 \quad \langle w, \nabla \quad {}^2\mathbf{f}(x) w \quad \overset{\lambda}{\rangle} + 0 \mathcal{L} \quad {}^2\mathbf{f}(x) = \mathbf{f} \quad$$

Now by strong convexity

$$\frac{1}{2} \quad \lambda^2 \quad \langle w, \nabla^2 \quad f(x) \ w \quad \rangle + 0 \ \dot{c} \quad {}^2) \quad \geq P\lambda \quad {}^2 \ \|w\|^2$$

$$= \sum \frac{1}{2} \quad \langle w, \nabla^2 f(x) w \rangle \quad + i \quad 0 \quad \frac{(\lambda^2)}{\lambda^2} \ge P \quad ||w||^2$$

Now as  $\lambda \downarrow 0$  (i.e;  $\lambda \rightarrow 0\dot{c}$  we have

$$\frac{1}{2} \quad \langle w, \nabla^2 \mathbf{f}(\mathbf{x}) \mathbf{w} \rangle \geq P \quad \|w\|^2$$

i.e; 
$$\langle w, \nabla^2 f(\mathbf{x}) w \rangle \ge 2P \|w\|^2$$

Thus  $\nabla^2$  f(x) is positive definite.

5. What are the major classes of convex optimization problems?

- a) Linear Programming problem
- b) Conic Programming problem
- c) Semi-definite Programming
- d) Quadratic convex programming under linear constraints
- e) Quadratic convex programming under quadratic constraint

• Linear Programming: min < ax >  
subject to  
$$Ax = b$$
  
 $x \ge 0$   
where C  $\in \mathbb{R}^n$ , A is a m  $\times$  n matrix, b  $\in \mathbb{R}$  m, &  $x \ge 0$   $\Rightarrow$   $x \in \mathbb{R}^n$ 

This is called linear programming in the standard form. Important feature: If a lower bound exists a minimizer exists.

• <u>Conic Programming</u> :

min < ax >

subject to  

$$Ax = b$$
  
 $x \in K$ 

where K is a pointed convex cone. The cone is called pointed if K  $\cap$  (-K) =  $\{0\}$ 

K for example could be the ice-cream cone or Lorenz-cone.

K= { x  $\in \mathbb{R}^n$  }:  $\sqrt{x_1^2 + x_2^2 + \dots + x_{n-1}^2} \dots \leq x_n; x_n \ge 0$  case the

above conic problem is called the second-order conic programming problem (SOCP for short).

Lorenz cone:



Lorenz cone is not a polyhedral cone.Semi- definite Programming :

 $S^{n} : \text{set of nXn systematic matrices} \\ \stackrel{+i}{S_{i}}^{n} : \text{set of nXn, systematic and positive semidefinite matrices} \\ S^{n}_{++} : \{X \quad \nabla \in S^{n}_{+} : X \text{ is positive definite}\} \\ \stackrel{+i}{S_{i}}^{n} \text{ is a convex cone but not polyhedral} \\ \text{Inner product in } S^{n} : iX, Y > i \text{ trace } (X, Y) \\ \min iC, X > i \\ iA : iX > i = b_{i} \\ X \quad \stackrel{+i}{\in S_{i}}^{n} \end{cases}$ 

• Semi definite programming or SDD for short is not a linear programming problem in matrices.

Quadratic convex programming with linear constraints.

$$\min \frac{1}{2} < x, Qx \qquad i \qquad + \qquad ic, x > i \qquad + \qquad d$$
  
subject to 
$$Ax = b$$
$$x \geq 0$$
$$Q \quad \in \qquad S^{n_{+}}, c \quad \in \mathbb{R}^{n} \quad , d \quad \in \qquad \mathbb{R}, A \text{ is a m} \qquad \times n \qquad \text{matrix and } x \geq 0 \Leftrightarrow x \in \mathbb{R}_{i}$$

<u>Important fact :</u> If a lower bound exists, then a minimizer exists. This is the celebrated Frank-Wolfe theorem.

• Quadratic convex programming with linear constraints

$$\min \frac{1}{2} \quad ix, Q_{0} \quad x > i \quad + \quad iC_{0}, \quad x > i \quad + \quad d_{0}$$
subject to
$$\frac{1}{2} \quad ix, Q_{i} \quad x > i \quad + \quad i \quad C_{i} \quad , \quad x > i \quad + \quad d_{i} \quad \leq 0$$

$$i = 1, \dots, m$$
where  $Q_{0}, \quad Q_{1}, \dots, \quad Q_{m}$  are positive semi-definite matrices,  $C_{0}, \quad C_{1}$ 

$$C_{m} \text{ are vectors in } R^{n} \quad \text{and } d_{0} \quad d_{1} \dots d_{m} \text{ are elements in R.}$$

## 6. What are saddle point conditions?

Consider the convex optimization problem (CP) min f(x) subject to  $g_i(x) \le 0$ , i-1,2,.....m

Construct the Lagrangian as follows

$$L(x, \lambda) = f(x) + \lambda_1 g_1(x) + \lambda_2 g_2(x) + \dots + \lambda_m g_m(x)$$

where  $\lambda = \begin{pmatrix} +i^m \\ \lambda 1 \\ \lambda m \end{pmatrix} \in R_i$  i.e;  $\lambda_i \ge 0$ , for all i=1,....m

A vector is (  $\dot{x}$ ,  $\dot{\lambda}$ )  $\overset{+\dot{\iota}^m}{\in R^n \times R_{\dot{\iota}}}$  is called a saddle point if

$$L\left(\begin{array}{c} \dot{x}, \lambda \end{array}\right) \leq \left(\begin{array}{c} \dot{x}, \dot{\lambda} \end{array}\right) \leq L\left(x, \begin{array}{c} \dot{\lambda} \end{array}\right), \text{ for all } x \in \mathbb{R}^{n} \text{ , and } \lambda \in \mathbb{R}^{n}_{d}$$

If solves convex optimization problem and slater condition holds, i.e; there exists  $\hat{x} \in \mathbb{R}^n$  s.t.

$$\begin{array}{ll} \mathbf{\dot{\iota}} g_i(\hat{x}) & \mathbf{\dot{\iota}} 0. \quad \forall \quad i=1,\dots,m \text{ then there exists } & \hat{\lambda} \in R_i \\ \text{i) } \mathsf{L} \left( \begin{array}{c} \hat{x}, \lambda \end{array} \right) \leq \left( \begin{array}{c} \hat{x}, \dot{\lambda} \end{array} \right) \leq \mathsf{L} \left( \mathbf{x}, \begin{array}{c} \dot{\lambda} \end{array} \right), \text{ for all } \mathbf{x} \in R^n \\ \text{ii) } & \hat{\lambda} g_i \end{array} ( \begin{array}{c} \hat{x} \end{array} ) = 0, \text{ } i=1,\dots,m \end{array}$$

If there exists a pair of  $(\dot{x}, \dot{\lambda}) \in \mathbb{R}^n \times R_i$  such that i) & ii) hold then  $\dot{x}$  solves (CP).